

AoPS Community

HMMT Invitational Competition 2017

www.artofproblemsolving.com/community/c437181 by CantonMathGuy

- 1 Kevin and Yang are playing a game. Yang has $2017 + \binom{2017}{2}$ cards with their front sides face down on the table. The cards are constructed as follows: For each $1 \le n \le 2017$, there is a blue card with n written on the back, and a fraction $\frac{a_n}{b_n}$ written on the front, where $gcd(a_n, b_n) = 1$ and $a_n, b_n > 0$. For each $1 \le i < j \le 2017$, there is a red card with (i, j) written on the back, and a fraction $\frac{a_i + a_j}{b_i + b_j}$ written on the front. It is given no two cards have equal fractions. In a turn Kevin can pick any two cards and Yang tells Kevin which card has the larger fraction on the front. Show that, in fewer than 10000 turns, Kevin can determine which red card has the largest fraction out of all of the red cards.
- **2** Let $S = \{1, 2, ..., n\}$ for some positive integer n, and let A be an n-by-n matrix having as entries only ones and zeroes. Define an infinite sequence $\{x_i\}_{i\geq 0}$ to be *strange* if: $x_i \in S$ for all i, - $a_{x_k x_{k+1}} = 1$ for all k, where a_{ij} denotes the element in the ith row and jth column of A. Prove that the set of strange sequences is empty if and only if A is nilpotent, i.e. $A^m = 0$ for some integer m.
- **3** Let v_1, v_2, \ldots, v_m be vectors in \mathbb{R}^n , such that each has a strictly positive first coordinate. Consider the following process. Start with the zero vector $w = (0, 0, \ldots, 0) \in \mathbb{R}^n$. Every round, choose an *i* such that $1 \le i \le m$ and $w \cdot v_i \le 0$, and then replace *w* with $w + v_i$.

Show that there exists a constant C such that regardless of your choice of i at each step, the process is guaranteed to terminate in (at most) C rounds. The constant C may depend on the vectors v_1, \ldots, v_m .

4 Let *G* be a weighted bipartite graph $A \cup B$, with |A| = |B| = n. In other words, each edge in the graph is assigned a positive integer value, called its *weight*. Also, define the weight of a perfect matching in *G* to be the sum of the weights of the edges in the matching.

Let G' be the graph with vertex set $A \cup B$, and (which) contains the edge e if and only if e is part of some minimum weight perfect matching in G.

Show that all perfect matchings in G' have the same weight.

5 Let S be the set $\{-1, 1\}^n$, that is, n-tuples such that each coordinate is either -1 or 1. For

 $s = (s_1, s_2, \dots, s_n), t = (t_1, t_2, \dots, t_n) \in \{-1, 1\}^n,$

define $s \odot t = (s_1 t_1, s_2 t_2, \dots, s_n t_n).$

Let c be a positive constant, let $f : S \to \{-1, 1\}$ be a function such that there are at least $(1-c) \cdot 2^{2n}$ pairs (s,t) with $s,t \in S$ such that $f(s \odot t) = f(s)f(t)$. Show that there exists a

AoPS Community

function f' such that $f'(s \odot t) = f'(s)f'(t)$ for all $s, t \in S$ and f(s) = f'(s) for at least $(1-10c) \cdot 2^n$ values of $s \in S$.

Act of Problem Solving is an ACS WASC Accredited School.